



Second Mid Semester Examination (March 15, 2019, Academic Session 2018-19)

IDC204 (Theory of Computation)

Maximum Marks: 20

Instructions

This booklet contains 8 printed pages, including this cover page, and 4 problems. Check if there are missing pages. Write your initials on the top of every page, in case the pages become separated. Attempt **ALL** problems. Read the problems carefully. Write all arguments precisely and do not leave **anything** to the instructor's¹ imagination. You need to sign the following **Academic Honour Code**, else your answers will not be evaluated.

"I affirm, on my honour, that I shall maintain my integrity and uphold the highest standards of academic conduct. I shall not receive or provide and shall not make any attempt to receive or provide an unlawful aid during this examination. I shall not tolerate any action that does not adhere to these words."

Signature *Abhigyan W. Medhi*

Name *Abhigyan W. Medhi*

Reg. No. *MS17108*

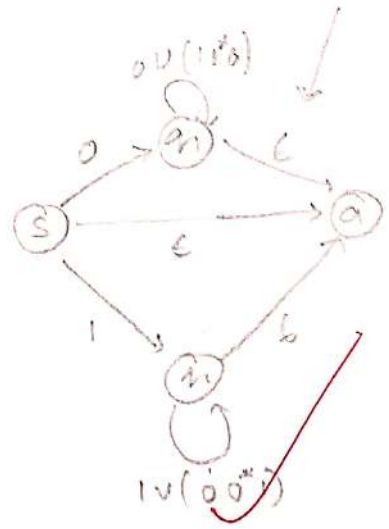
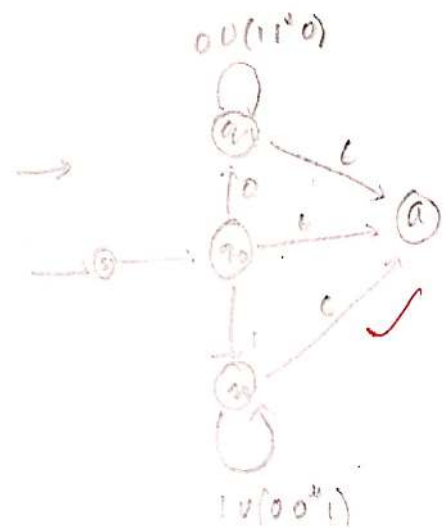
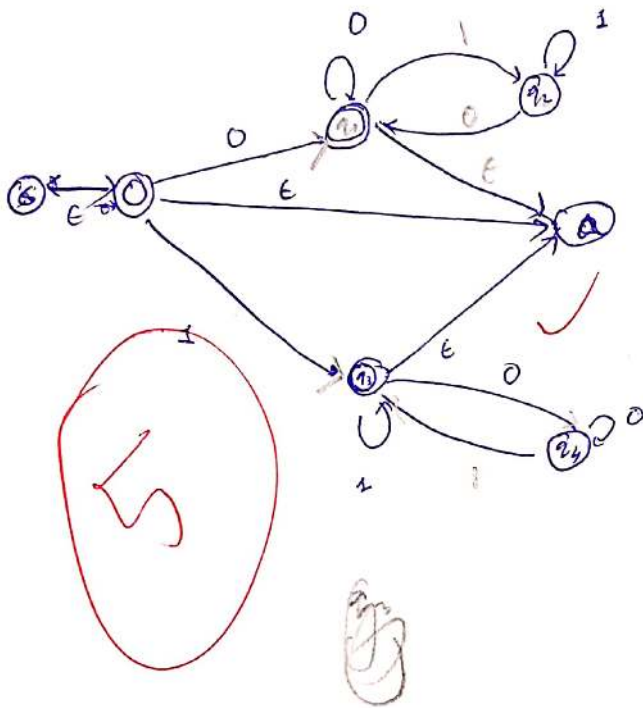
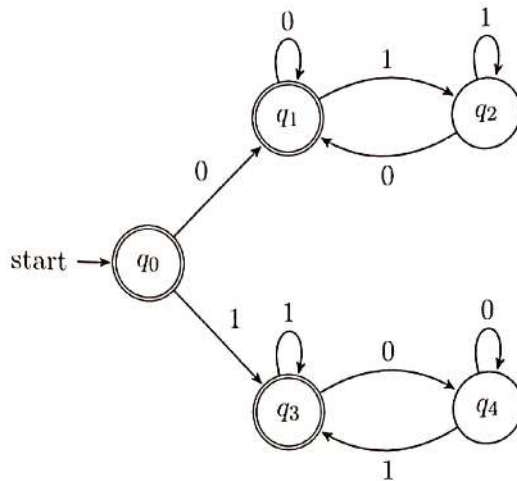
The following rules apply:

- **Organize your work**, in a reasonably neat and coherent way, in the space provided. Work scattered all over the page without a clear ordering will receive very little credit. Illegible work will not receive credit.
- **Mysterious or unsupported answers will not receive credit.** A correct answer, unsupported by calculations, explanation, or algebraic work will receive no credit; an incorrect answer supported by substantially correct calculations and explanations *might* still receive partial credit.
- **Do your rough work in the blank pages attached at the end of this booklet.** No other space will be provided for the rough work.

| Problem | Points | Score |
|---------|--------|-------|
| 1 | 5 | 5 |
| 2 | 5 | 5 |
| 3 | 5 | 5 |
| 4 | 5 | 5 |
| Total: | 20 | 20 |

¹Instructor: Amit Kulshrestha

1. (5 points) Find a regular expression that evaluates to the language recognized by the following DFA.



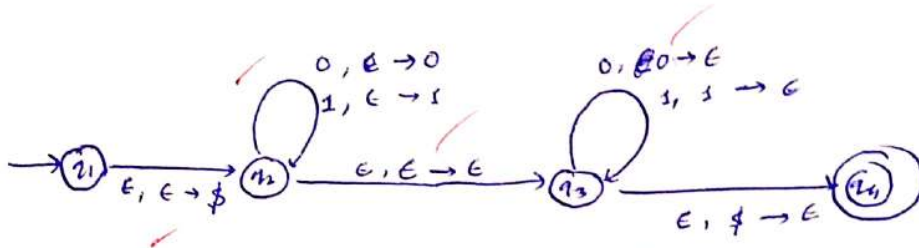
$$S \xrightarrow{\epsilon \cup (0(0^*(1^*0)^*))} A$$

$$\cup (1(1^*(0^*1)^*))$$

$$S \xrightarrow{\epsilon \cup (0(0^*(1^*0)^*)) \cup (1(1^*(0^*1)^*))} A$$

2. (5 points) Construct a pushdown automaton that recognizes palindromes of even length, i.e. the language $\{ww^R : w \in \{0,1\}^+\}$. Here w^R denotes the reverse of the string w .

$$PDA = \{ Q, \Sigma, \Gamma, \delta, q_1, q_4 \}$$



$$Q = \{ q_1, q_2, q_3, q_4 \}$$

$$\Sigma = \{0, 1\}$$

$$\Gamma = \{0, 1, \$\}$$

$$\delta(q, p, \gamma) = \begin{cases} (q_1, \epsilon, \epsilon) \rightarrow (q_2, \$) & \gamma \in \Gamma \\ (q_2, 0, \epsilon) \rightarrow (q_2, 0) & p \in \Sigma \\ (q_2, 1, \epsilon) \rightarrow (q_2, 1) & p \in \Sigma \\ (q_2, \epsilon, \epsilon) \rightarrow (q_3, \epsilon) & \gamma \in \Gamma \\ (q_3, 0, 0) \rightarrow (q_3, \epsilon) & \gamma \in \Gamma \\ (q_3, 1, 1) \rightarrow (q_3, \epsilon) & \gamma \in \Gamma \\ (q_3, \epsilon, \$) \rightarrow (q_4, \epsilon) & \gamma \in \Gamma \end{cases}$$

3. (5 points) Consider the language $A = \{0^n 1^n 2^n : n \geq 0, n \in \mathbb{N}\} \subseteq \{0, 1, 2\}^*$. Determine if A is regular.

3/ Let $s \in A$ be a regular language
and p be its pumping length

$$\text{Let } s = 0^p 1^p 2^p \quad |s| > p$$

$\therefore \exists xyz$ s.t. $xy^i z \in A$ for all i

Case I y has only 0's

Then $xy^i z \notin A$ as no. of 1's will be greater

Case II y has only 1's

Then $xy^i z \notin A$ as no. of 0's will be greater

Case III ~~Then $xy^i z \in A$~~ y has only 2's

Then also, $xy^i z \notin A$ as no. of 2's will be greater

Case IV y has both 1's and 2's

Then $xy^i z \notin A$ as order will not be maintained

Similarly, y cannot have both 2's & 0's and 0's & 1's

Case V y has all 1's, 2's and 3's

$xy^i z \notin A$ as the order will not be maintained

$\therefore A$ is not a regular language.

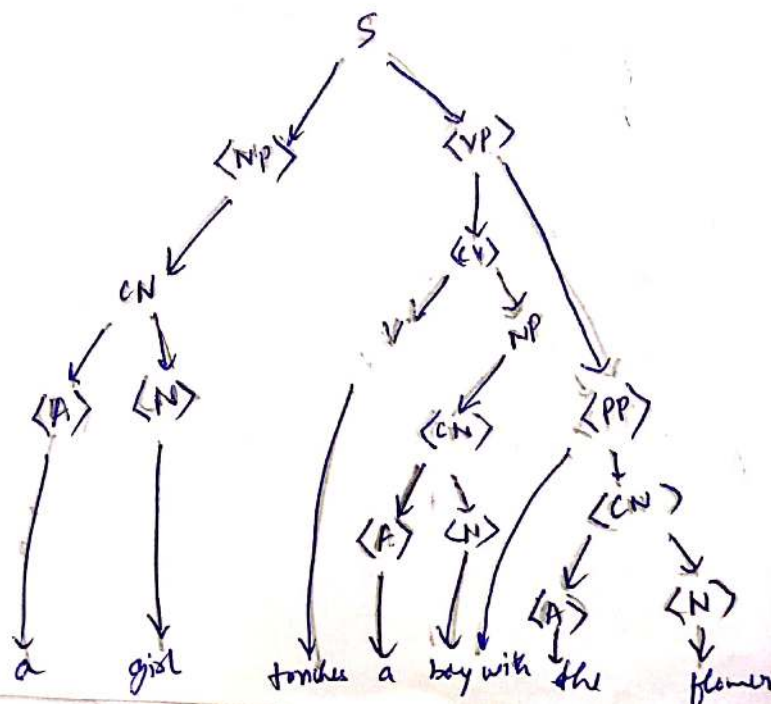
4. (5 points) Consider the following grammar.

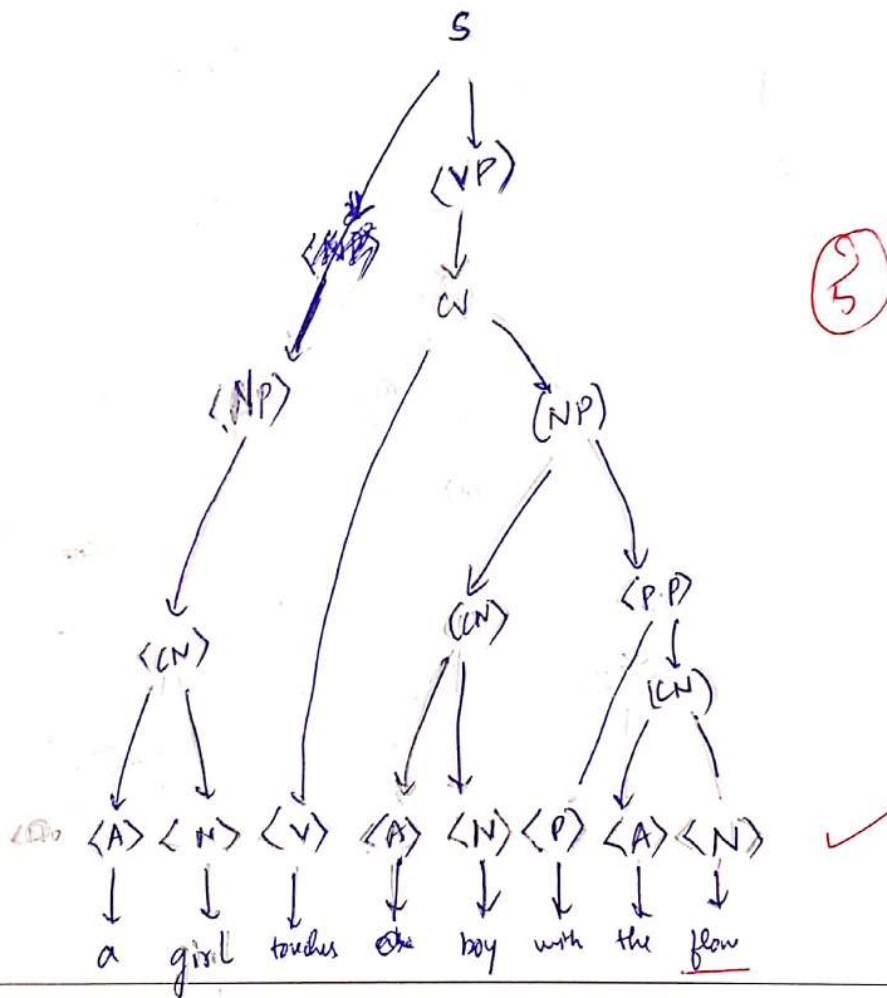
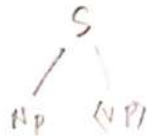
- $\langle S \rangle \rightarrow \langle NP \rangle \langle VP \rangle$
- $\langle NOUN-PHRASE \rangle \rightarrow \langle NOUN-PHRASE \rangle \langle VERB-PHRASE \rangle$
- $\langle NOUN-PHRASE \rangle \rightarrow \langle COMPLEX-NOUN \rangle \mid \langle COMPLEX-NOUN \rangle \langle PREPOSITION-PHRASE \rangle$
- $\langle VERB-PHRASE \rangle \rightarrow \langle COMPLEX-VERB \rangle \mid \langle COMPLEX-VERB \rangle \langle PREPOSITION-PHRASE \rangle$
- $\langle PREPOSITION-PHRASE \rangle \rightarrow \langle PREPOSITION \rangle \langle COMPLEX-NOUN \rangle$
- $\langle COMPLEX-NOUN \rangle \rightarrow \langle ARTICLE \rangle \langle NOUN \rangle$
- $\langle COMPLEX-VERB \rangle \rightarrow \langle VERB \rangle \mid \langle VERB \rangle \langle NOUN-PHRASE \rangle$
- $\langle ARTICLE \rangle \rightarrow a \mid the$
- $\langle NOUN \rangle \rightarrow boy \mid girl \mid Mohali \mid flower \mid spectacles \mid orange$
- $\langle VERB \rangle \rightarrow touches \mid eats \mid sees \mid talks \mid likes \mid reads$
- $\langle PREPOSITION \rangle \rightarrow on \mid under \mid with \mid after$

Show that the string $w = "a\ girl\ touches\ a\ boy\ with\ the\ flower"$ is ambiguous by exhibiting two leftmost derivations of w . Draw parse trees for the two derivations.

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$S \Rightarrow \langle NP \rangle \langle VP \rangle \Rightarrow \langle A \rangle \langle N \rangle \langle V \rangle \langle NP \rangle \Rightarrow \langle A \rangle \langle N \rangle \langle V \rangle \langle NP \rangle$
 $\Rightarrow \langle A \rangle \langle N \rangle \langle V \rangle \langle CN \rangle \langle PP \rangle \Rightarrow a \langle N \rangle \langle V \rangle \langle A \rangle \langle N \rangle \langle PP \rangle$
 $\Rightarrow a\ girl \langle V \rangle \langle A \rangle \langle N \rangle \langle PP \rangle \Rightarrow a\ girl\ touches \langle A \rangle \langle N \rangle \langle PP \rangle$
 $\Rightarrow a\ girl\ touches\ a \langle N \rangle \langle PP \rangle \Rightarrow a\ girl\ touches\ a\ boy \langle PP \rangle$
 $\Rightarrow a\ girl\ touches\ a\ boy\ with\ the\ flower.$





This is the last printed page of this booklet. Any work done beyond this sheet will be considered rough.

